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Canberra

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I further certify that the above application is now proceeding in the name of VIRTUAL BUSINESS ASSOCIATES PTY LIMITED pursuant to the provisions of Section 113 of the Patents Act 1990.

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Tenth day of January 2000

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MANAGER PATENT OPERATIONS

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S & F Ref: 438805

ORIGINAL

AUSTRALIA

Patents Act 1990

PROVISIONAL SPECIFICATION FOR THE INVENTION ENTITLED:

Secure Multi-Point Data Transfer System

Name and Address of Applicant: Virtual Business Associates Pty Limited ~~Lyal Sidney Collins, of 1/37 Walton Crescent, Abbotsford, New South Wales, 2046, AUSTRALIA~~ 27 Werona Avenue, Killara NSW 2071.

Name of Inventor(s): Lyal Sidney Collins

This invention is best described in the following statement:

MESSAGE IDENTIFICATION WITH CONFIDENTIALITY, INTEGRITY, AND SOURCE AUTHENTICATION

Field of the Invention

The present invention relates to the encoding and transmission of secure messages, in particular relating to aspects of confidentiality, integrity and auditability of messages in terms of authentication and integrity checking. In addition, the invention relates to reliable operation of such messaging functions in a network environment in which transmission delay and lost or duplication of messages can occur.

Background of the Invention

The advent of secure storage and processing devices such as smart-cards, coupled with the increasing use of practicable electronic commerce technology, has highlighted shortcomings in secure message transfer technology. This relates in particular to the robustness and auditability of secure messages when transmitted over different types of "best effort" networks.

Fundamental requirements for electronic commerce include the ability to transmit and receive messages with an acceptable level of confidentiality and integrity, where this level depends on the particular commercial application. In addition, reliable authentication of these messages, namely identification and verification of the source of a received message is also needed to ensure that fraudulent transactions are not being initiated.

Emerging best effort networks such as wireless and the Internet, place additional demands on messaging technology, since message delay, loss and occasionally duplication does occur.

Standard cryptographic and authentication functions often exact a commercially prohibitive penalty on secure messaging, because of their requirement for significant overhead data and associated complex equipment to provide the cryptographic and/or

authentication functions. Available techniques have also not been proven to be reliable or efficient in the context of the aforementioned best effort networks.

It is an object of the present invention to ameliorate one or more disadvantages of the prior art.

Summary of the Invention

According to a first aspect of the invention, there is provided a method for encoding and transmitting by an originating device of a secure message the method comprising the steps of;

(a) generating by a first process using an application identifier and an application value of a message value;

(b) combining the message value with one or more first secret values, said secret values being known substantially only to the originating device and one or more intended recipient devices of the message, to establish a secret message value;

(c) applying the secret message value and the message to an encoding process to form a secure message block; and

(d) combining an address with a device identifier, the application identifier, the application value and the secure message block, to form a secure message for transmission which is decodable by the one or more of said intended recipient devices which thereby recover the message, the address , the device identifier, the application identifier and the application value.

According to another aspect of the invention, there is provided a method for reception of a securely transmitted message by a recipient device the method comprising the steps of:

(i) extracting an application identifier and application value from a received secure message;

(j) generating by a first process using the application identifier and application value of a message value;

(k) extracting a device identifier from the secure message, whereby one or more secret values known substantially only to an originating device and the one or more intended recipient devices of the message are generated according to a second process using the device identifier and the application identifier;

(l) combining the message value with the one or more secret values, to establish a secret message value;

(m) extracting a secure message block from the secure message; and

(n) applying the secret message value and the secure message block to a decoding process to form the securely transmitted message, this message having been securely transmitted by the originating device.

Brief Description of the Drawings

A number of embodiments of the invention are described with reference to the drawings, in which:

Fig. 1 depicts secure communication between Issuers and device-holders;

Fig. 2 depicts the sourcing of devices and device applications from different issuers;

Fig. 3 depicts a device holder performing authentication in relation to a device;

Fig. 4 illustrates incorporation of secret values into Issuer and device-holder devices;

Fig. 5 illustrates a preferred embodiment for producing a secret message unique value ;

Fig. 6 depicts a preferred embodiment for production of a transmission data block;

Fig. 6a depicts an embodiment for production of a transmission data block with confidentiality and integrity protection;

Fig. 6b depicts another embodiment for production of a transmission data block with confidentiality and integrity protection;

Fig. 7 depicts another embodiment for production of a transmission data block;

Fig. 8 illustrates a preferred embodiment for reception of the secret message unique value ; and

Fig. 9 illustrates a decoding process for recovery of the message.

Detailed Description

The term "comprising" as used herein has the inclusive meaning of "having" or "including" and not the exclusive meaning of "consisting only of".

The term "unique" is used herein in one of two ways. In the first instance, it is used as a label e.g. "Unique Application Value". In the second instance, it is used to indicate the manner of parameter value selection for a number of parameters. For example, "secret values are preferably unique values" is taken to mean that secret values are chosen in a manner as to minimise the likelihood that two secret values will have the same value.

Electronic network communications involve both originators of messages, and recipients of those messages. Some communication systems dealing with applications like e-mail handling, financial services, and directed research information acquisition involve a large number of individuals communicating uni-directionally and/or bi-directionally with a small number of servers or hosts. Systems of this type are characterised by communication paths which are "many to one" or "many to few".

Turning to Fig. 1, an Issuer 100 communicates with a number of Device-holders 104 and 106 across a network 108. Another Issuer 102 communicates with the same device-holders 104 and 106, and with other device-holders (not shown) across the network 108.

Fig. 2 shows how the communication referred to in relation to Fig. 1 is performed by the Issuer 100 (see Fig. 1) using an Issuer device 200 to communicate with the device-holder 104 by means of the device-holder device 202. The Issuer device 200 communicates across the network 108 to the device holder 202 using corresponding applications 206 and 208 respectively which are incorporated into the respective devices 200 and 202. The Issuer device 200, the device-holder device 202,

and the applications 206 and 208 are either proprietary or commercial products, and are generally available from different suppliers in the market. This requires that the applications 206, 208 and devices 200, 202 comply with appropriate interface and interworking standards. In the rest of the description, communication between issuer and device-holder and communication between issuer device and device-holder device are taken to have the same meaning unless a contrary intention is stated.

The Issuer device 200 and the device-holder device 202 ensure the confidentiality and integrity of communication, independent of the type of network infrastructure 108. They provide confidentiality and message integrity even in the event that messages are delayed, corrupted, or delivered in a different sequence to the one in which they were transmitted.

The Issuer device 200 communicates with device-holder device 202 for a variety of different purposes. These purposes include administrative functions such as exchanging logon ID/passwords and exchanging account information. They also include sending, and receiving electronic mail, sending and receiving purchase information in relation to a purchase, or transacting purchases. Each communication type is associated with a particular application in the Issuer device 200 and a corresponding application in the device-holder device 202. A suite of applications (e.g. 214 and 206) in the Issuer device 200 can be supplied as an integrated set of applications, or alternatively as modular software applications from different sources. The same applies in regard to a suite of applications in the device-holder device 202.

Fig. 3 illustrates how a device-holder 104 can in some circumstances, typically at the issuer's discretion, be required to perform an authentication procedure, as depicted by arrow 302, in regard to the device-holder device 202. This authentication procedure 302 can, for example, take the form of an exchange of password identification, or can use a biometric identification procedure such as placing the device-holders thumb on a special purpose thumb-print sensor. Alternatively, passive authentication can be achieved by mere possession of the device-holder device 202.

Where required by the particular application (e.g. exemplified by 206, 208), the aforementioned authentication procedure provides authentication information which can be incorporated into the communication messages. For example, communications dealing with requests for health, financial or computer system access information commonly require, as a prerequisite to answering the request, a reliable indication that the information request has originated from a device and/or application which is known to, and authorised by, the information provider. Furthermore, the information provider must be sure that the device making the request is being used by a user who is in turn authorised to make such a request. In this case, the authentication information can be incorporated into each message, to enable the message recipient to assess the authentication status of a message at the time of receipt. The authentication or message identification information can be used for network performance assessment, in order to estimate the integrity and efficiency of the communication system, and the individual communication links. In addition, the authentication information can be used as a basis for establishing the origin, destination, sequence and timing of messages. This is usable, for example, in customer dispute resolution situations, as substantiating evidence.

The aggregate level of security provided by the Issuer device 200, the application (e.g 206 and 208), and the device-holder device 202 is specified by the Issuer 100, to comply with his requirements and those of the device-holders 104 and 106. The Issuer will normally specify a required level of security based upon risk management assessment of the Issuer's business requirements. Tamper-resistant card-reading terminals and smart-cards are an example of a particular issuer device 200 and associated device-holder device 202 respectively in the case, for example, where the Issuer is a bank, and the device-holder is a bank customer.

The Issuer device 200 and the device-holder device 202 (see Fig. 2) are generally arranged to erase sensitive data values held in storage if the devices are subjected to tampering or damage. Typically, in the case of multiple applications 214 and 206 being resident in the Issuer device 200 or device holder device 202, an

operating system within the issuer device 200 provides secure access control to data on a per-application basis. The level of security associated with inter-application access varies with the type of messaging application, for example, financial or health applications being more security-intensive than lower priority e-mail massaging.

Having regard to Fig. 4, the Issuer device 200 is able to store secret values 400 in a secure manner. The secret values 400 will typically be at least 64 bits long, but preferably will be 112 bits or greater in length (i.e. the length of a double key according to the digital encryption standard (DES), or other symmetric encryption process such as LOKI, IDEA, RC4 and so on). The secret values 400 being such length preclude practicable brute force attacks which could otherwise be feasibly used to deduce the secret values 400.

The Issuer device 200 and the device-holder device 202 are arranged to allow one or more secret values 400 known only to the Issuer's device 200 and the device-holders device 202 to be stored in both the Issuer device 200 and the device-holder device 202. Typically, two unique secret values 400 will be used, one for message origination, and the second for message reception. Other situations or applications however, only require a single secret value 400. An example of this is an application for secure identification, encryption and decryption of data or files for backup or external storage purposes, where a single device acts as both the originator 200 and recipient 202 at differing times.

Provision of distinct secret values 400 for each application (e.g. 206, 208) within a device (e.g 200, 202) provides for reliable and single valued indication of both the device and application that originate a particular message. The Issuer device 200 and the device-holder device 202 are engineered in a fashion as to preclude misuse of secret values 400.

The secret values 400 are preferably unique values. This ensures not only that particular applications have different secret values 400, but also that any secret value 400 has a low probability of being the same as secret values 400 used in any other device holder 202 or application e.g 218.

The corresponding applications 206 and 208 are assigned application identity values 406 and 414, to permit identification of an application or purpose for a particular message. This identification can vary between applications, or between versions of the same application. The application identity (406) can be either a numeric value (e.g "1, 2, 3, 4, 5, 6"), or a more descriptive text string (e.g. "ABC banking system", or, "ABC banking system logon step 1").

Each device-holder device 202 and issuer device 200 is allocated a device identifier 408,416 which might, for example, be a serial number. This provides a unique identifier for each device. The device identifier 408, 416 allows the issuer device to know which device-holder device originates a message.

The issuer device 200 maintains, in some secure fashion, a record of the device identifier 408, the relevant application identifier 414, and the secret values 400 associated with all the devices e.g. 202 and/or applications e.g. 208 issued by the Issuer. The Issuer device 200 stores multiple secret value sets, each set being specific to both a device and an application, while each application within a device will contain a secret value set. The Issuer stores information regarding both the devices which are registered to communicate with it, and the applications which the registered devices contain.

This is exemplified in the following table, which illustrates typical data maintained by the issuer device 200, illustrating how a number of different secret values SV^s, SV^r, SV^i can be associated with a record set.

DID	**Application ID (AID)**	**Secret Value Send(SV^s)**	**Secret Value Receive (SV^r)**	**Secret Value Integrity (SV^i)**
123653	remote access v1.01	247EB4BC8EF52	2F667C42C2C02	
123654	remote access v1.01	10A6B1C8ED9F9	48009F1CCE203	
	1098756	99A73E7D456A8		

DID	Application ID (AID)	Secret Value Send(SV^s)	Secret Value Receive (SV^r)	Secret Value Integrity (SV^i)
123655	ABC savings account Cash Management v2.9	3C768B8A71C31 2906F8812A346 C459EAC53F55	4789239EFAAB1 387FEA1B4755C4 7E89564CA2313	2906F8812A34E C459EAC53F5A3
123656	ABC savings account	83E76FC890323	345F7898AC1F5	11FF045A67897

Table 1.

Devices can contain multiple applications, which communicate with this issuer. Thus device 123654 contains a first application entitled "Remote Access v1.01" and another application entitled "1098756".

Fig. 5 below illustrates how the secret value or, in the case shown in Table 1 the secret values, SV^k are combined with the application identifier e.g. 406, 414, the device identifier e.g. 408, 416 and a message related value e.g. 412.

A single instance of the application 206 within the device 200 can require one or more secret values. Thus with reference to Table 1. application "Remote Access v1.01" requires a secret value SV^s whose value is "10A6B1C8ED9F9" for ensuring confidentiality in the send direction. The same application further requires a secret value SV^r whose value is "48009F1CCe203" for ensuring confidentiality in the receive direction.

Devices associate corresponding details on applications, secret values and those Issuers with which the device has been registered. Extracting the DID and AID fields from a received message enables the Issuer to retrieve the appropriate secret value(s). A device retrieves appropriate secret value(s) by virtue of the Issuer's DID and AID fields within a received message.

The application identifier 406 permits a message-originating device to tag a specific message with the identifier 406 when delivering it to a recipient device.

For auditing and indexing purposes an application-unique value 412 is assigned to each message transmitted. This application-unique value 412, when combined with the device identifier 416 and the application identifier 406, permits reliable indexing of every message within a system or network. This indexing is related to the message, the device, and the application. The application-unique value 412 can be a simple counter within the application 206 or the Issuer device 200. Alternatively, time and/or date information or a combination of the aforementioned parameters can be used. The range of the application-unique value 412 encompasses the expected working life (i.e. the total expected number of messages sent/received during the lifetime) of the device (e.g. 200) and the application (e.g. 206). A binary value of 32 bits or 10 decimal digits normally suffices for this purpose.

Fig. 5 illustrates a preferred embodiment of the message origination process. The Application identifier 406 and the unique application value 412 are combined in a process 500 to create a message unique value 502. The combination process 500 produces a message unique value 502 which is individual to the specific input combination of the application identifier 406 and the unique application value 412. Cryptographic techniques such as symmetric encryption, using Cipher Block Chaining (BCB) or another cipher feedback mode, keyed hash functions, or hash functions such as SHA-1 and MD5 fulfil this required functionality. In contrast, exclusive OR (XOR) functions are generally not suitable, since the resulting message unique value 502 will not be unique. If a keyed function such as the symmetric key encryption based one way function is used, using the unique application value 412 as the key value will marginally increase the work factor for some forms of attack.

The message unique value 502 is combined with the secret value 400 in combination process 504 to form a secret message unique value 506. The secret message unique value 506 is substantially unique to the particular message, device and application. It is noted that the secret value 400 is logically associated with the device identifier 416 and application identifier 406.

The combination process 504 can be implemented using the symmetric encryption based one way functions used in the financial industry, and/or hash functions such as SHA-1 and MD5. The use of non-reversible combination processes 504 is preferred to encryption processes, in order to isolate the secret value 400 from possible recovery due to brute force attacks, should one or more secret message unique values 506 be compromised in any manner.

Turning to Fig. 6, the secret message unique value 506 is combined with message data 600 in an encoding process 602. This process 602 can be selected appropriately to provide symmetric key encryption for confidentiality, or for providing a message integrity mechanism, such as a Message Authentication Code (MAC) or keyed hash function, or simply as a secret one-time value for use within a higher level protocol. More details on MACs can be found in Australian Standard 2805 and in ANSI X9 Standards and similar documents.

The encoding process 602 outputs a secure message block 604 which is unique to the message 600, device 200 and application 206. This encoding process 602 binds the device identifier 416, the application identifier 406, the application unique value 412, and the secret values 400 to the message 600.

Message data or content is formatted according to the needs of the issuer and device holder. Message length and/or content can be arbitrarily arranged. Encryption and/or message integrity functions are incorporated together with the message data as exemplified by a transmission data block 606. The transmission data block 606 takes the form of three major components, namely the secure message block 604, control data 610, and addressing data 612. The control data 610 consists of the device identifier 416, the application identifier 406, and the unique application value 412. The addressing data 612 consists of a destination address 618, a source address 616, and optionally, ancillary data 614. The format of the transmission data block 616 is determined by the Issuer 100 (see Fig. 1).

Considering Fig. 6 with reference to Fig. 1, the secure message block 604 is opaque, that is indecipherable, to all network entities apart from the intended recipient e.g. 104.

The format and arrangement of the addressing data 612 is related to network functionality and not directly to the messaging functions of authentication and integrity assurance. Addressing data 612 is thus specific to the purpose, network and processing devices being employed by the Issuer device 200 and device-holder device 202.

This arrangement also allows the same device identifier 408 to be used at multiple network addresses 618, 616. Alternatively, redundant issuer devices each with a distinct device identifier can be accessed at the same network address.

Fig. 6a depicts a situation where both confidentiality and integrity protection are required. In a first embodiment, two encoding processes 602 and 603 are applied in parallel, process 602 for confidentiality and process 603 for integrity. Two distinct secret values SV^c (for confidentiality) and SV^i (for integrity) are used to produce two secret message unique values 632 and 630 respectively. These are applied to the corresponding processes 602 and 603 together with message data 600 to produce two secret message blocks 620 and 604 respectively. The transmission data block 622 is then constructed to contain the two secret message blocks 604 and 620. Symmetric key encryption can be used for confidentiality, and Message Authentication Code (MAC) or keyed hash function can be used for integrity.

In a second embodiment, still having regard to Fig. 6a, if both confidentiality and integrity are required, the first secret value SV^c is used to produce the secret message value 632 using process 504 (see Fig. 5). The secret message value 632 is then combined with message data 600 in confidentiality encoding process 602 to produce the secure message block 620 and thereafter, a transmission data block. The second secret value SV^i is then used to produce the secret message value 630 using process 504 (see Fig. 5). Thereafter, the secret message value 630 is encoded in integrity encoding process 603 together with the aforementioned transmission data block to produce the secure message block 604. This is then used to form a

transmission data block which has been iteratively encoded to provide both confidentiality and integrity protection.

Turning to Fig. 6b, in a third embodiment where both confidentiality and integrity are required , the message data 600 is combined with the secret message value 506 in the confidentiality encoding process 602 to form a confidentiality secure message block 604. The same secret message value 506 is in parallel combined with a MAC Variant 1000 in XOR process 1002 to output an integrity secret message value 1008. This secret message value 1008 is then combined with the message data 600 in the integrity encoding process 1004 to form an integrity secure message block 1006. The confidentiality secure message block 604 and the integrity secure message block 1006 are then incorporated into transmission data block 606. MAC Variants are described in AS2805, ANSIX9, and similar standards.

Where both confidentiality and integrity protection are required, the sequence of processing may be decided according to the needs of the issuer. Thus, processing for confidentiality protection may be applied prior to processing relating to integrity protection, or alternatively, the processing may be performed in the reverse order.

Fig. 7 illustrates another embodiment whereby the secret message unique value 506 is combined with message data 600 and the message unique value 502 in encoding process 602 to produce the secure message block 700 and thereafter to form transmission data block 702. This enables the message recipient to detect whether the incoming transmission data block 702 has been altered or corrupted during transmission, without performing a complete message reception procedure, and also allows utilisation of partially intact messages.

Fig. 8 illustrates a preferred embodiment which relates to decoding of the transmission data block 606. The application unique value 412 and application identifier 406 are extracted from the incoming transmission data block 606, and combined in the process 500 to recreate the message unique value 502. The combination process 500 is the identical process used in the message transmission process as described in Fig. 5.

The device identifier 408 and the application identifier 406 are extracted from the transmission data block 606 and used to retrieve the appropriate secret value 400 by means of a secret value retrieval process 802.

The recreated message unique value 502 is combined with the retrieved secret value 400 in the combination process 504, in order to derive the secret message unique value 506. The combination process 504 is identical to the process utilised to combine the message unique value 502 and the secret value 400 in the transmission process described in Fig. 5.

Turning to Fig. 9, the secret message unique value 506 is utilised to decode the secure message block 604 in a decoding process 900, in order to produce the original message data 600. The decoding process 900 is the inverse process to the encoding process 602 (see Fig. 6). Thus if the encoding process 602 implemented symmetric key encryption, i.e. related to confidentiality, then the decoding process 900 decrypts the secure message block 604 using the unique value 506. If the encoding process 602 (see Fig. 6) implemented a message integrity mechanism such as a MAC or keyed hash function, then the decoding process 900 verifies the integrity of the secret message block 604 against message corruption or tampering, using MAC or keyed hash techniques, or both, as applicable.

Where the message unique value 502 is included with message data 600 in forming the secure message block 700 (see Fig. 7), application of the secret message unique value 506 to the secure message block 604 which contains the transmitted message unique value 502 in decoding process 900 allows detection of errors in the transmission data block 606 if it contains errors in the control data 610 (see Fig. 6) and parts of the secure message block 604.

Thus the message recipient device 202 and application (e.g. 208) utilise publicly disclosed items of information transmitted within the transmission data block 606 and one or more shared secret values 400 to uniquely identify the contents of the transmission data block 606.

Any other receiving entity with access to the network 108 and having authorised access to appropriate secret values 400 or secret message unique value 506 can also identify a corresponding transmission data block 606, and the incorporated destination device and/or application for purposes of metering, charging, quality control or law enforcement purposes. Where only the secret message unique value 506 has been provided for these purposes, prior and subsequent messages which use the secret value 400 are not compromised.

The foregoing describes only some embodiments of the present invention, and modifications obvious to those skilled in the art, can be made thereto without departing from the scope of the invention.

ASPECTS OF THE INVENTION

The following numbered paragraphs set forth aspects of the invention:

1. A method for encoding and transmitting by an originating device of a secure message the method comprising the steps of;

(a) generating by a first process using an application identifier and an application value of a message value;

(b) combining the message value with one or more first secret values, said secret values being known substantially only to the originating device and one or more intended recipient devices of the message, to establish a secret message value;

(c) applying the secret message value and the message to an encoding process to form a secure message block; and

(d) combining an address with a device identifier, the application identifier, the application value and the secure message block, to form a secure message for transmission which is decodable by the one or more of said intended recipient devices which thereby recover the message, the address , the device identifier, the application identifier and the application value.

2. The method according to paragraph 1 whereby an association of the device identifier, the application identifier, and the application value substantially uniquely identifies the originating device and a purpose of the message and/or the application, and a identifier for the message, such message identification being bound with the message content by virtue of the encoding process.

3. The method according to paragraph 1 whereby the encoding process in step (c) comprises either:

(e) a symmetric encryption process, or

(f) an integrity process using keyed hash or symmetric encryption techniques, or

(g) a process including both symmetric encryption and keyed integrity, or

(h) including the secret message value in a higher level messaging protocol.

4. A method for reception of a securely transmitted message by a recipient device the method comprising the steps of:

(i) extracting an application identifier and application value from a received secure message;

(j) generating by a first process using the application identifier and application value of a message value;

(k) extracting a device identifier from the secure message, whereby one or more secret values known substantially only to an originating device and the one or more intended recipient devices of the message are generated according to a second process using the device identifier and the application identifier;

(l) combining the message value with the one or more secret values, to establish a secret message value;

(m) extracting a secure message block from the secure message; and

(n) applying the secret message value and the secure message block to a decoding process to form the securely transmitted message, this message having been securely transmitted by the originating device.

DATED this FOURTH DECEMBER 1998
~~Lyal Sidney Collins~~
Virtual Business Associates Pty Limited
Patent Attorneys for the Applicant/Nominated Person
SPRUSON & FERGUSON

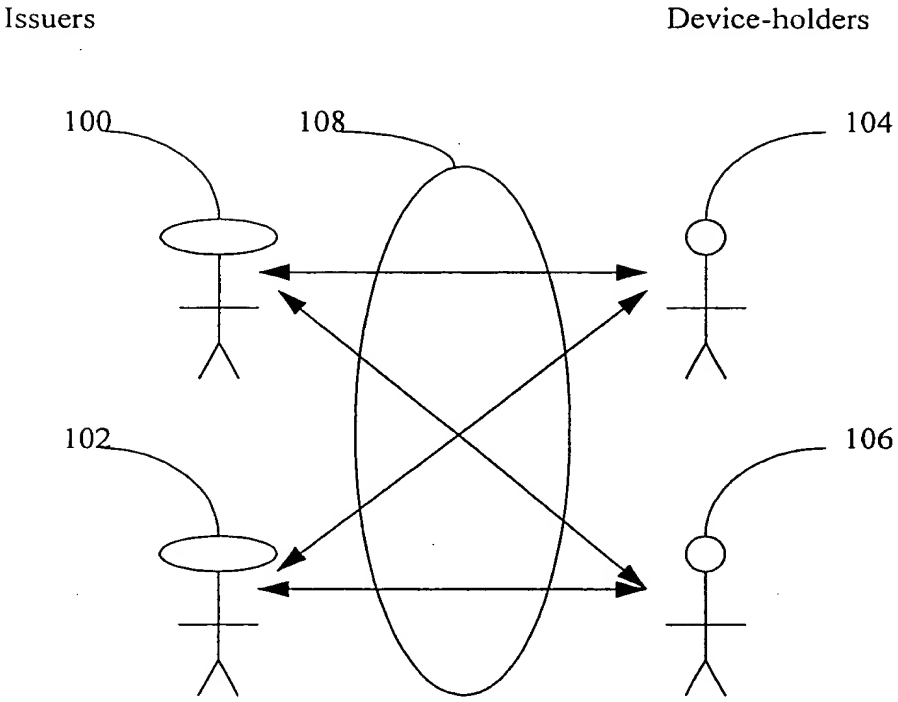
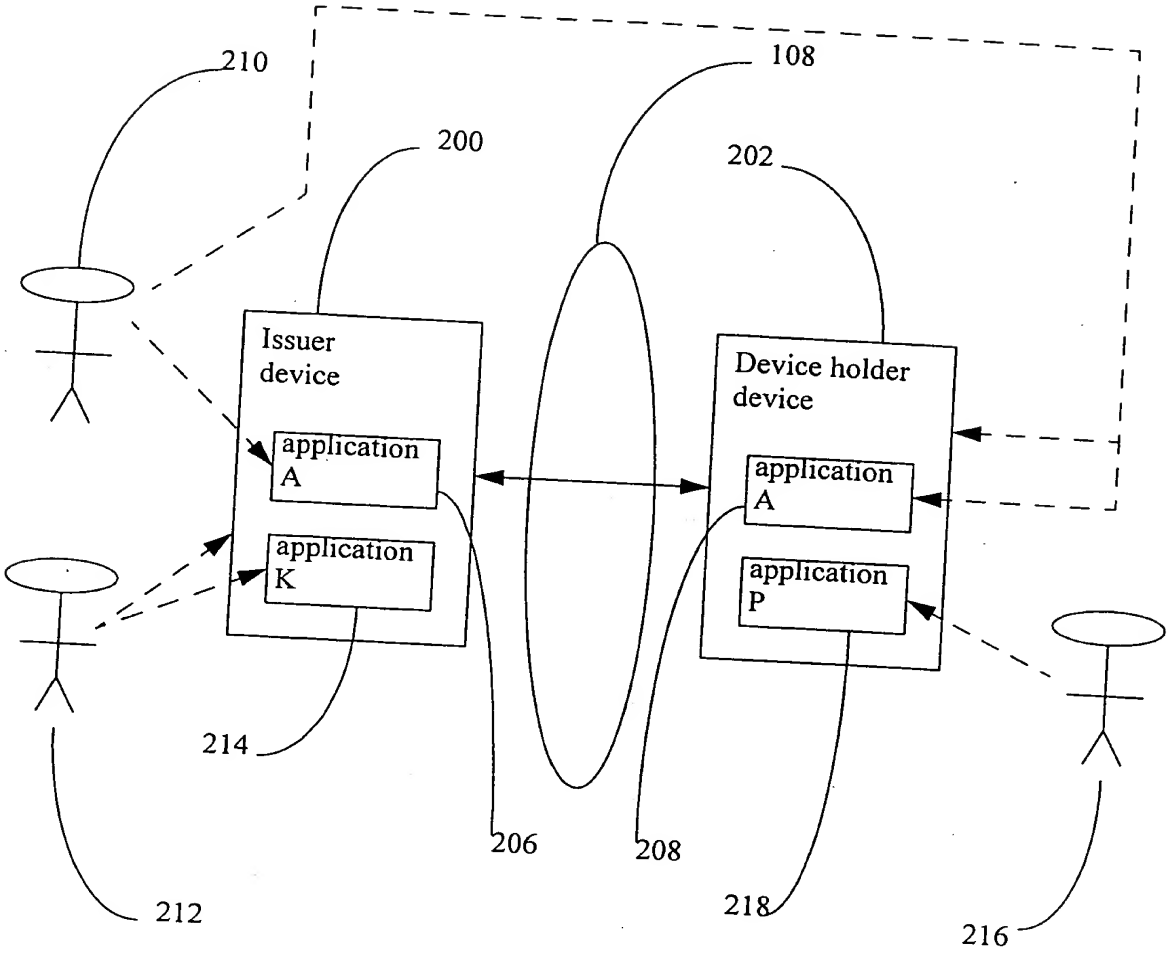


Fig. 1


210
108
200
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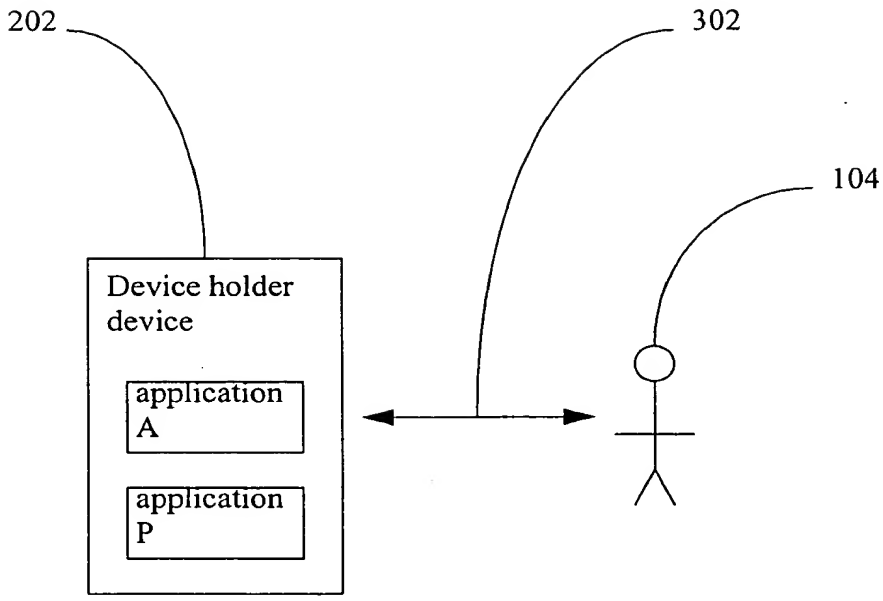
Issuer device
application A
application K
Device holder device
application A
application P
214
206
208
218
212
216

Fig. 2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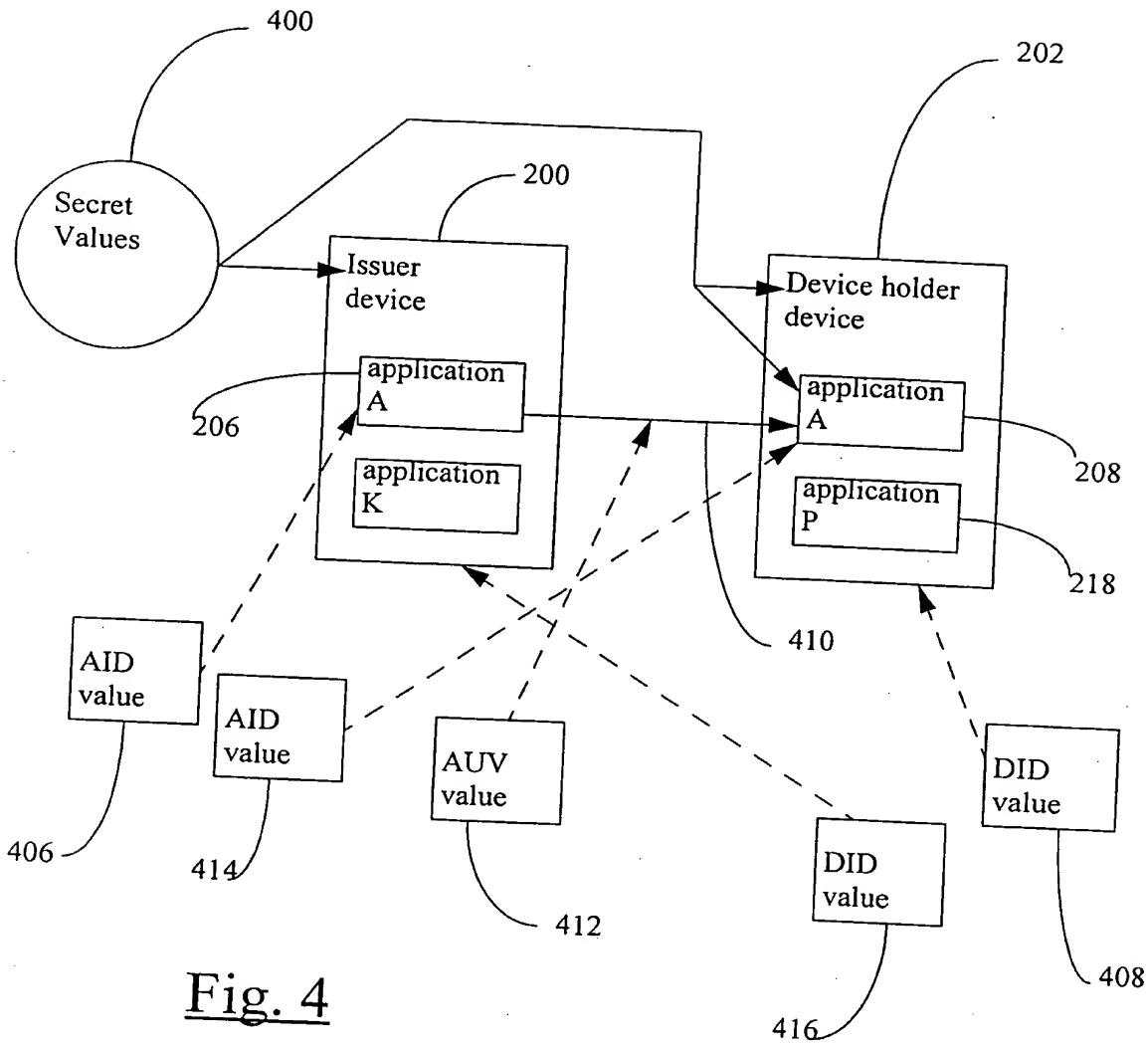
302
104
Device holder device
application A
application P

Fig. 3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Secret Values
200
Issuer device
application A
application K
206
202
Device holder device
application A
application P
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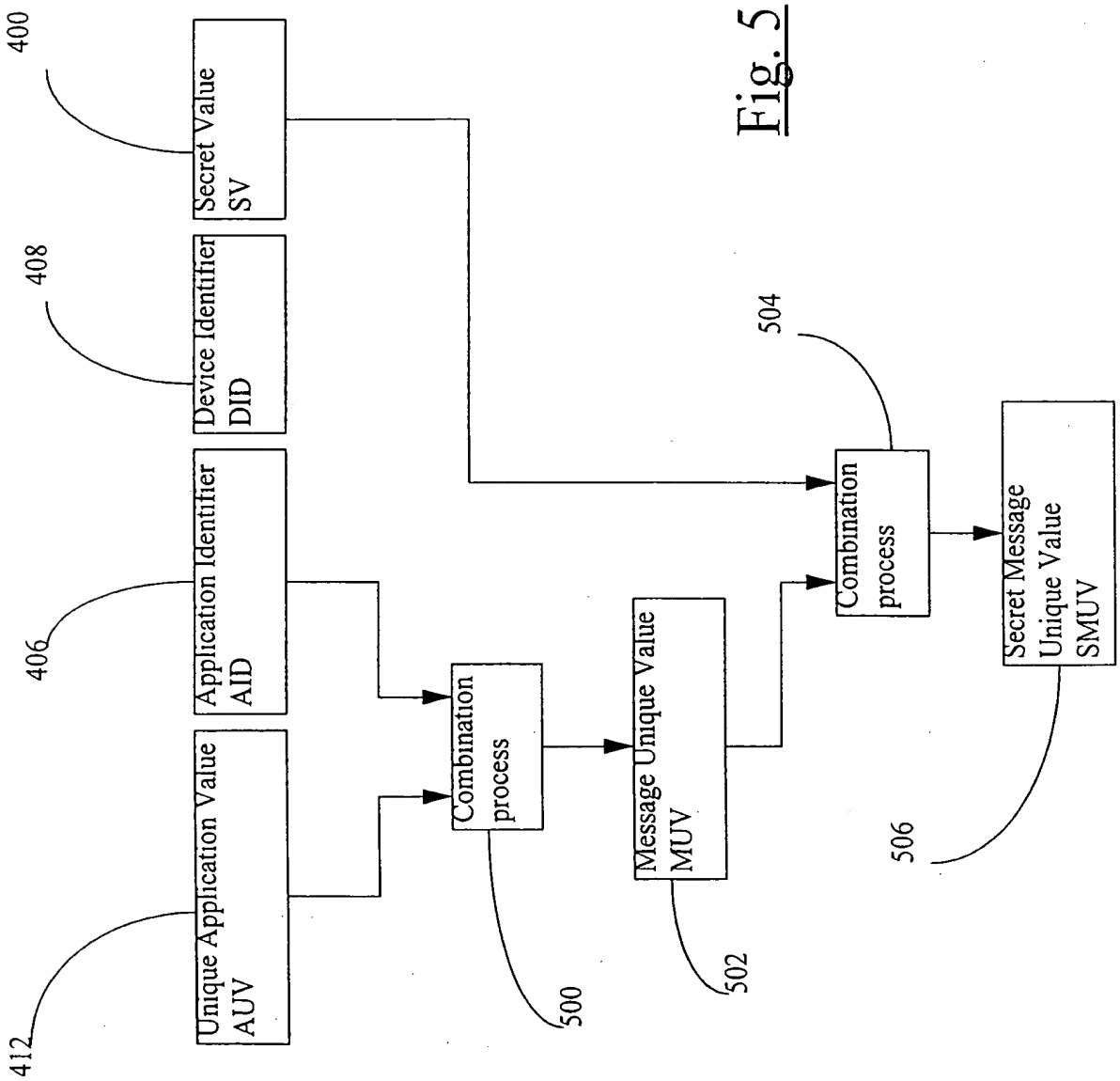
DID value
DID value
406
414
412
416
408

Fig. 4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DID
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AID
412
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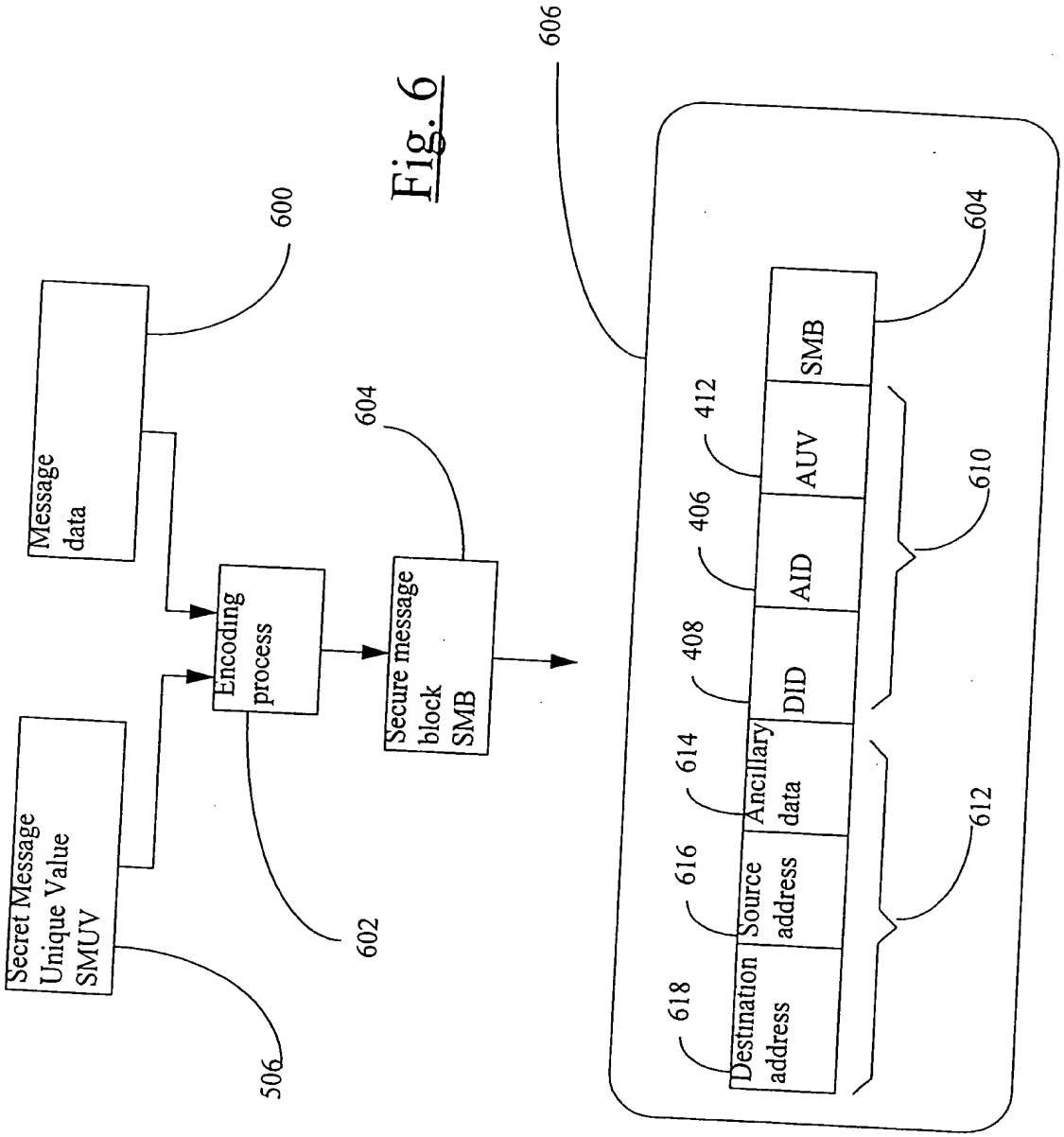
Fig. 5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Secure message block SMB
604
606
SMB
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Source address
Destination address
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412
406
408
614
616
618
610
612

Fig. 6

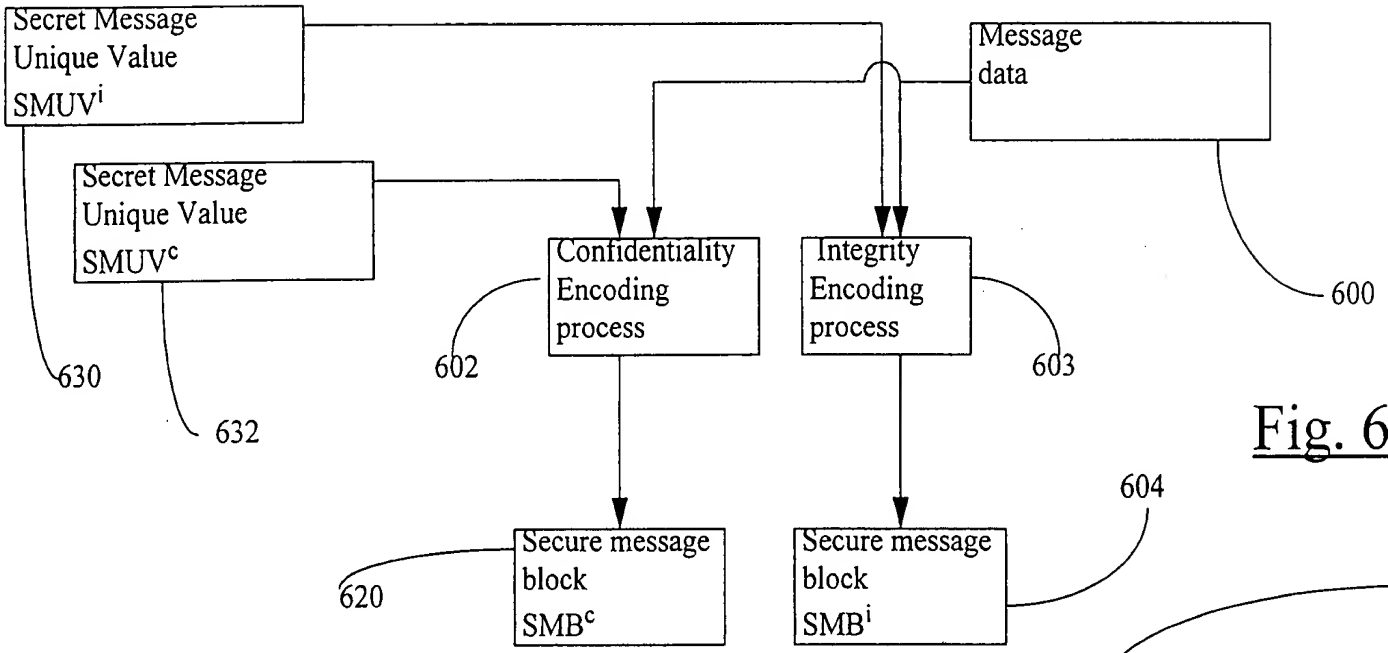
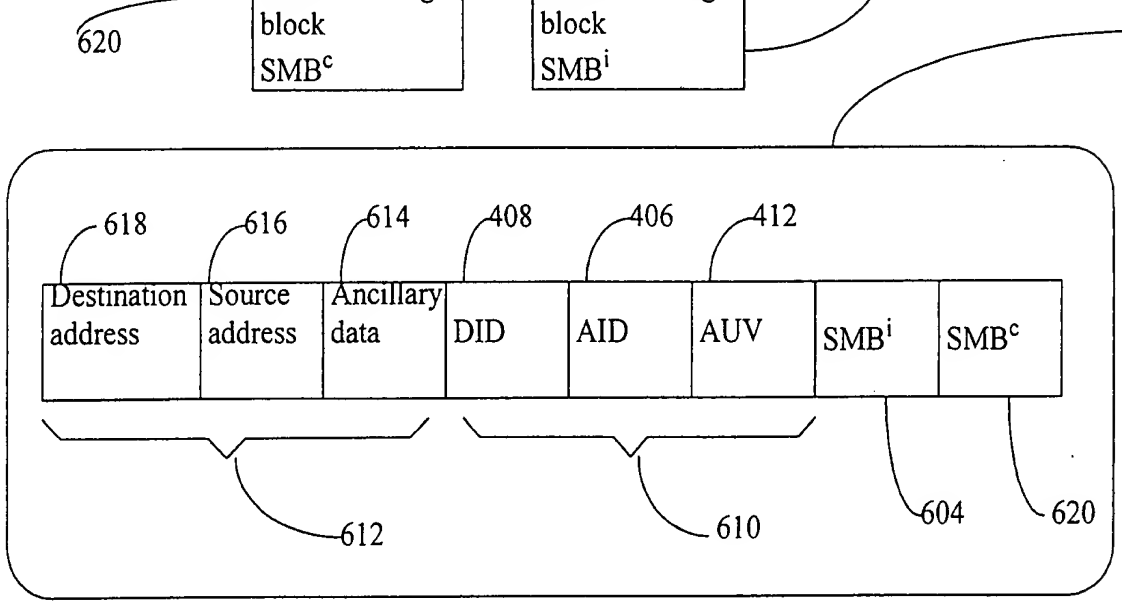

Secret Message Unique Value SMUV^i
Secret Message Unique Value SMUV^c
Message data
Confidentiality Encoding process
Integrity Encoding process
630
632
602
603
600
620
Secure message block SMB^c
Secure message block SMB^i
604

Fig. 6a


622
618
616
614
408
406
412
Destination address
Source address
Ancillary data
DID
AID
AUV
SMB^i
SMB^c
612
610
604
620

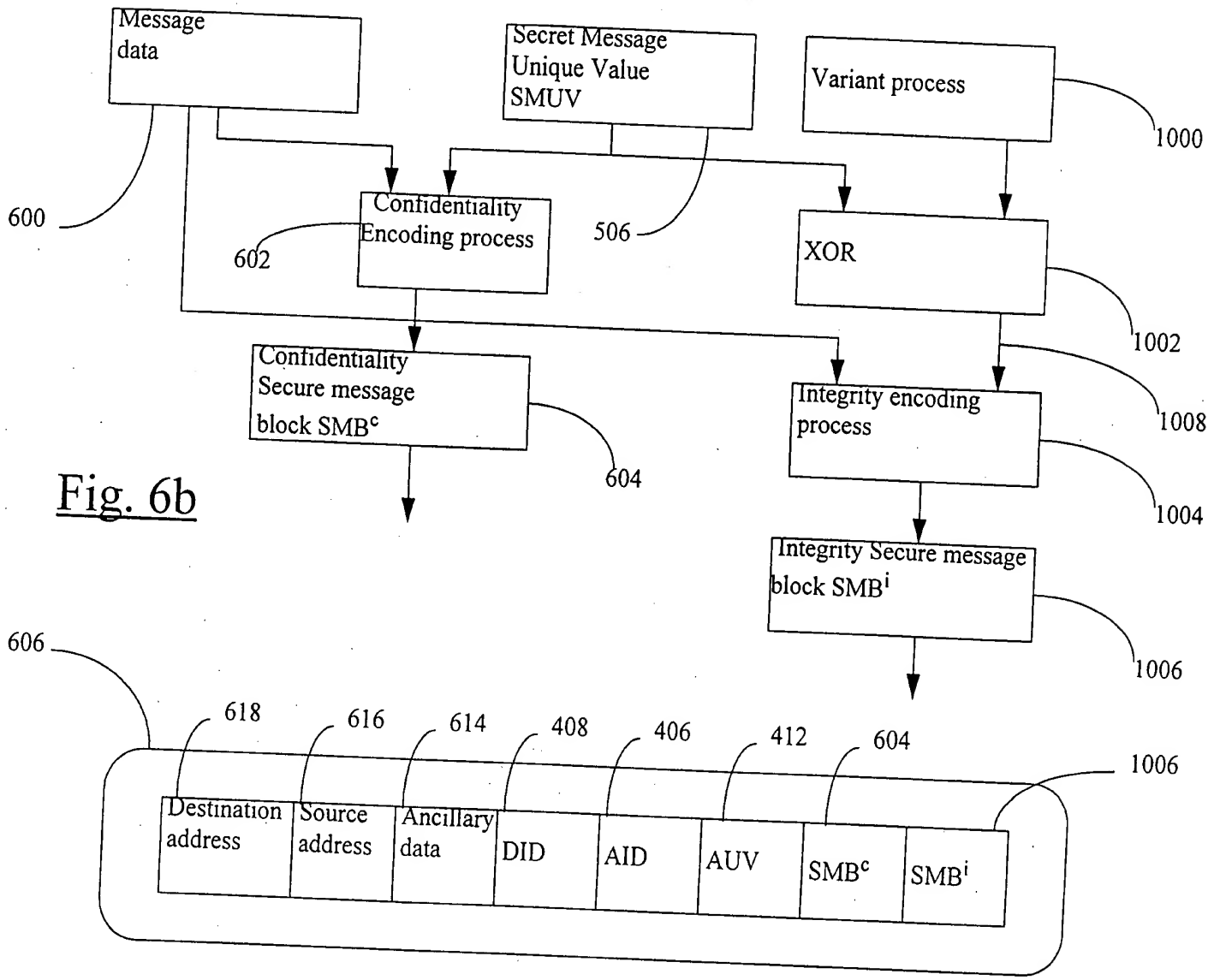

Message data
Secret Message Unique Value SMUV
Variant process
1000
600
Confidentiality Encoding process
602
506
XOR
1002
Confidentiality Secure message block SMB^c
604
Integrity encoding process
1008
1004
Integrity Secure message block SMB^i
1006
606
618
616
614
408
406
412
604
1006
Destination address
Source address
Ancillary data
DID
AID
AUV
SMB^c
SMB^i

Fig. 6b

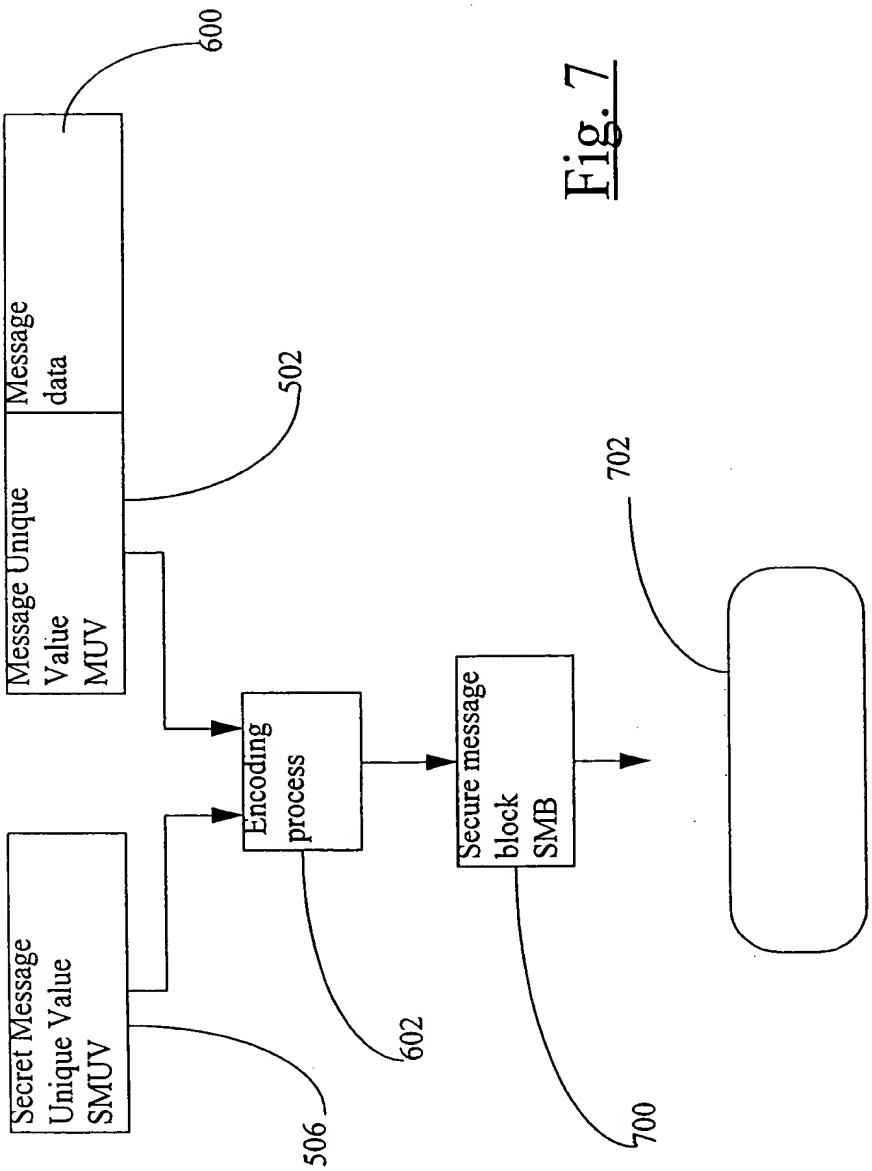

600
Message Unique Value MUV
Message data
502
Secret Message Unique Value SMUV
506
Encoding process
602
Secure message block SMB
700
702

Fig. 7

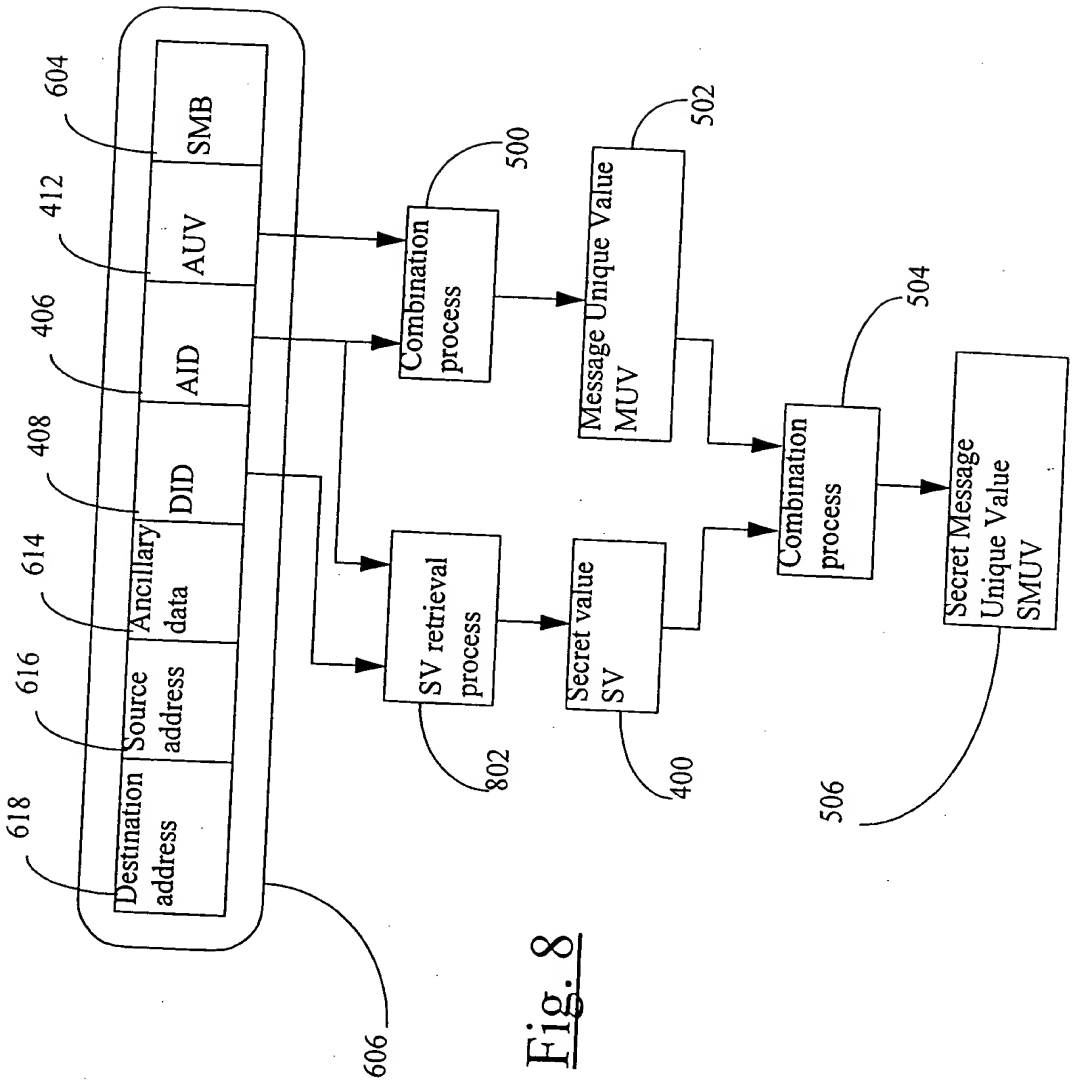

Destination address
Source address
Ancillary data
DID
AID
AUV
SMB
606
618
616
614
408
406
412
604
SV retrieval process
802
Combination process
500
Secret value SV
400
Message Unique Value MUV
502
Combination process
504
Secret Message Unique Value SMUV
506

Fig. 8

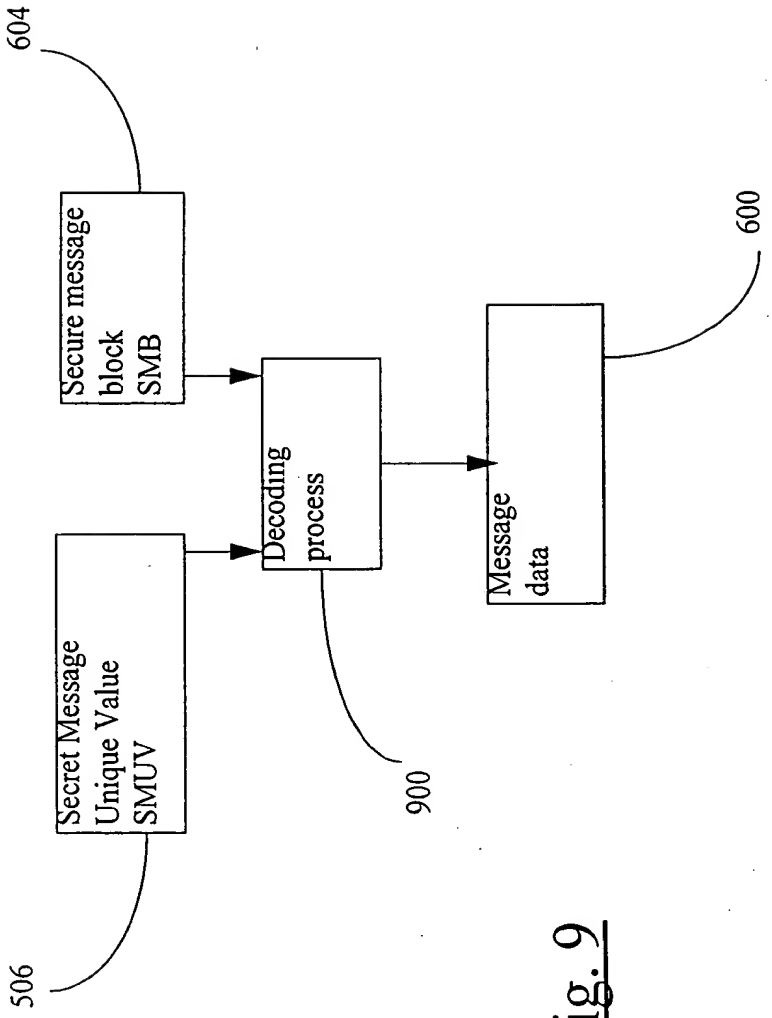

604
Secure message block SMB
Secret Message Unique Value SMUV
506
Decoding process
900
Message data
600

Fig. 9